

Binary Division

$$321 \div 7$$

DMSB

↓
1^u
6^s
↓
5000000

$$\begin{array}{r} 45 \\ 7 \overline{) 321} \\ \underline{32} \\ 0 \\ \underline{08} \\ 41 \end{array}$$

$$321 \div 7 = 45 \text{ r } 6.$$

$$1101101 \div 101 = 0010101 \text{ r } 100$$

$$\begin{array}{r} 0010101 \\ 101 \overline{) 1101101} \\ \underline{110} \\ 001 \\ \underline{001} \\ 000 \\ \underline{000} \\ 000 \end{array}$$

Only on unsigned numbers; resign if necessary.

$$101101 \div 11 = 001111$$

$$\begin{array}{r}
 001111 \\
 11 \overline{) 101101} \\
 \underline{101} \\
 00 \\
 \underline{00} \\
 00 \\
 \underline{00} \\
 0
 \end{array}$$

$$45 \div 3 = 15 \checkmark$$

Fractional components

$$11.011 \div 11$$

$$\begin{array}{r}
 01.001 \\
 11 \overline{) 11.011} \\
 \underline{11} \\
 011 \\
 \underline{011} \\
 0
 \end{array}$$

$$3.375 \div 3 = 1.125$$

$$1 \div 11$$

$$\begin{array}{r}
 0.010101 \\
 11 \overline{) 1.000000} \\
 \underline{10} \\
 00 \\
 \underline{00} \\
 00 \\
 \underline{00} \\
 00 \\
 \underline{00} \\
 00 \\
 \underline{00} \\
 00 \\
 \underline{00} \\
 00
 \end{array}$$

etc.....

Floating point numbers

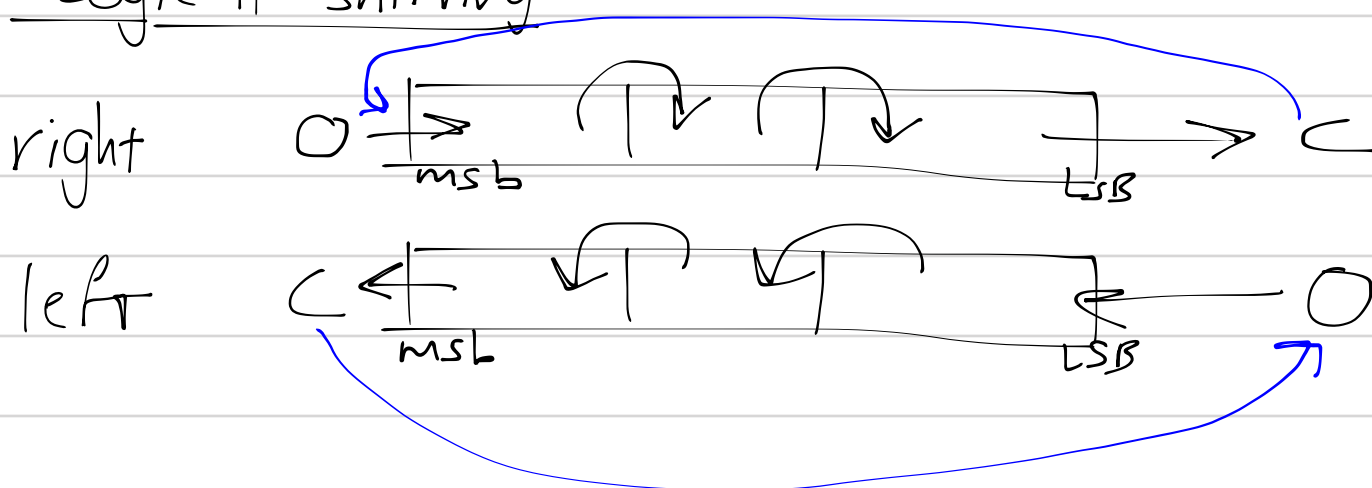
$$y = M \times b^E$$

M is the mantissa. in Q0.m format ALWAYS.
b is the base.

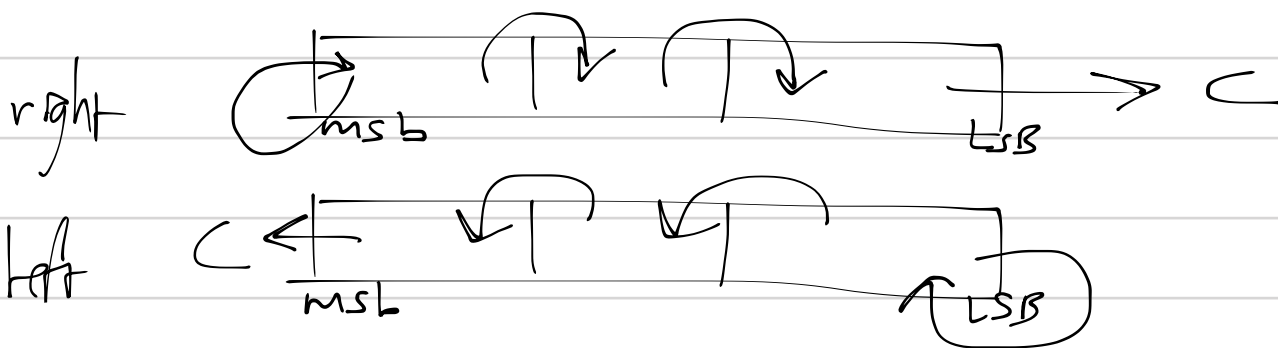
E is the exponent in Qn.0 format ALWAYS

y has format Qmen
Q_{2e3}

Logical shifting



Arithmetic shifting



$$10110 \text{ ras } 1 = 11011 [0]$$

$$-16 + 4 + 2 = -10$$

$$-16 + 8 + 2 + 1 = -5$$

$$11011.0 \text{ ras } 1 = 11101.1$$

$$= -2.5$$

$$01101 \text{ ffs } 1 = 011010 \approx \times 2.$$

Normalisation

Normalisation is the method by which precision is maximised.

$$\begin{array}{l} 0.11010e0000 \rightarrow -5+0 \Rightarrow \text{LSB order} \\ 0.01101e0001 \rightarrow -5+1 \Rightarrow \text{LSB order} \end{array}$$

$$(0 + 0.5 + 0.25 + 0.0625) \times 2^0 = 0.8125$$

$$(0 + 0.25 + 0.125 + 0.03125) \times 2^1 = 0.8125$$

$$(35.82 \times 10^3) \div 100 = (35.82 \times 10^{3-2})$$

$$\begin{array}{l} 1.1101e0111 \\ = 1.1010e0110 \quad (\text{left logical shift } 1) \\ = 1.0100e0101 \quad (\text{left logical shift } 2) \\ = 0.1000e0100? \quad \text{overflow. } \times \end{array}$$

To normalise: logical left shifts until $\text{MSB} = \overline{\text{MSB}_1}$
for every shift subtract 1 from exponent.

$$\begin{array}{l} 0.00101e0111 \\ 0.10100e0101 \quad \& \text{ left shifts.} \end{array}$$

$$\begin{array}{l} \rightarrow 2^{-5} \times 2^7 = 2^2 \\ \rightarrow 2^{-5} \times 2^5 = 2^0 \end{array}$$

Justification

Justification is the process by which exponents are matched

NB: Always done on normalised numbers.

Right arithmetic shift until smaller exponent matches bigger one.

$$\text{Justify } 0.110e001 \text{ on } 1.011e011$$

$$= 0.011e010$$

$$= 0.001e011$$

NB information lost
Present in guard bit

Floating point addition

$$0.01101e0110 + 1.1101e0010$$

1. Normalise (both numbers)
2. Justify on largest exponent
3. Do addition
4. Normalise answer.

$$= 0.11010e0101 + 1.01000e0000 \text{ (normalised)}$$

$$= 0.11010e0101 + 1.1111e0101$$

$$\begin{array}{r} 011010 \\ + 111111 \\ \hline 1011001 \end{array}$$

$$\begin{array}{r} 011010 \\ + 111111 \\ \hline 1011001 \end{array}$$

normal signed addition.

carry, no overflow

$$= \underline{0.11001e0101 \text{ (normalised)}} \rightarrow$$