ELEN 4017

Network Fundamentals Lecture 8

Purpose of lecture

Chapter2: Application Layer

• Principles of network applications



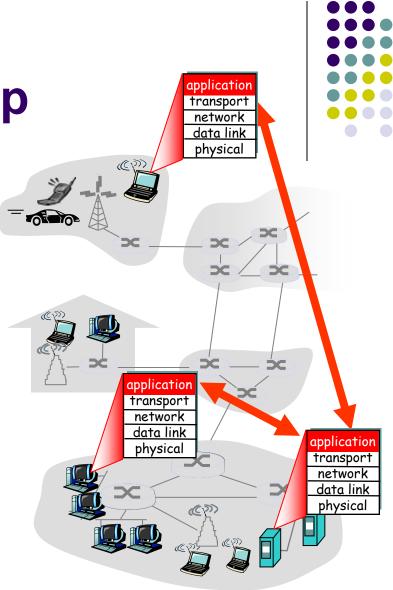
Creating a network app

write programs that

- run on (different) *end systems*
- communicate over network
- e.g., web server software communicates with browser software

No need to write software for network-core devices

- Network-core devices do not run user applications
- applications on end systems allows for rapid app development, propagation

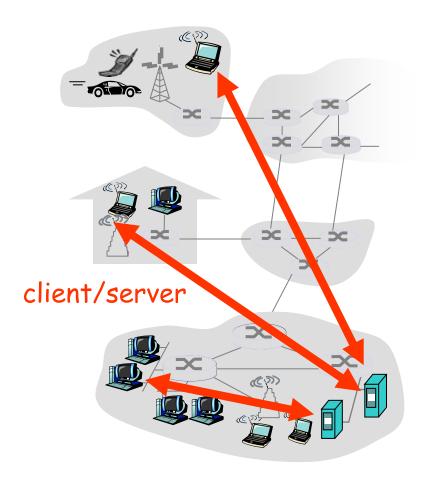




Application architectures

- Client-server
- Peer-to-peer (P2P)
- Hybrid of client-server and P2P

Client-server architecture



server:

- always-on host
- permanent IP address
- server farms for scaling

clients:

- communicate with server
- may be intermittently connected
- may have dynamic IP addresses
- do not communicate directly with each other

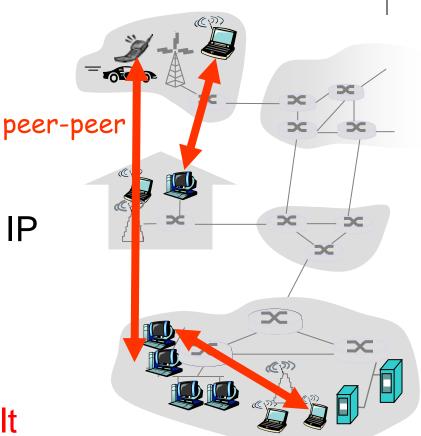




Pure P2P architecture

- no always-on server
- arbitrary end systems directly communicate
- peers are intermittently connected and change IP addresses

Highly scalable but difficult to manage



Hybrid of client-server and P2P



Skype

- voice-over-IP P2P application
- centralized server: finding address of remote party:
- client-client connection: direct (not through server)

Instant messaging

- chatting between two users is P2P
- centralized service: client presence detection/location
 - user registers its IP address with central server when it comes online
 - user contacts central server to find IP addresses of buddies



Processes communicating

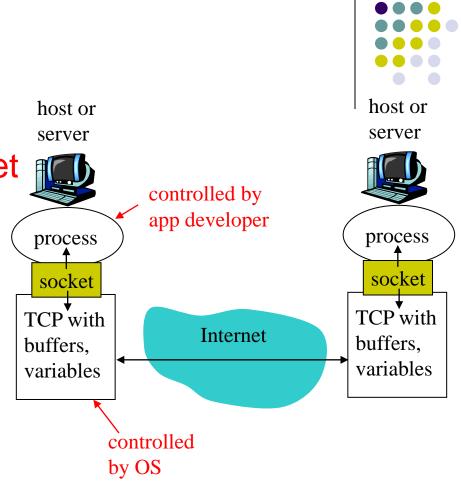
- Process: program running within a host.
- within same host, two processes communicate using inter-process communication (defined by OS).
- processes in different hosts communicate by exchanging messages

Client process: process that initiates communication Server process: process that waits to be contacted

 Note: applications with P2P architectures have client processes & server processes

Sockets

- process sends/receives messages to/from its socket
- socket analogous to door
 - sending process shoves message out door
 - sending process relies on transport infrastructure on other side of door which brings message to socket at receiving process



• API: (1) choice of transport protocol; (2) ability to fix a few parameters (lots more on this later)

Addressing processes

- to receive messages, process must have identifier
- host device has unique 32-bit IP address
- <u>Q</u>: does IP address of host suffice for identifying the process?



Addressing processes

- to receive messages, process must have identifier
- host device has unique 32-bit IP address
- Q: does IP address of host on which process runs suffice for identifying the process?
 - <u>A</u>: No, *many* processes can be running on same host

- identifier includes both IP
 address and port numbers associated with process on host.
- Example port numbers:
 - HTTP server: 80
 - Mail server: 25
- to send HTTP message to gaia.cs.umass.edu web server:
 - IP address: 128.119.245.12
 - Port number: 80
- more shortly...

App-layer protocol defines



- Types of messages exchanged,
 - e.g., request, response
- Message syntax:
 - what fields in messages & how fields are delineated
- Message semantics
 - meaning of information in fields
- Rules for when and how processes send & respond to messages

Public-domain protocols:

- defined in RFCs
- allows for interoperability
- e.g., HTTP, SMTP
- Proprietary protocols:
- e.g., Skype



RFC793 – TCP

September 1981

Transmission Control Protocol

PREFACE

This document describes the DoD Standard Transmission Control Protocol (TCP). There have been nine earlier editions of the ARPA TCP specification on which this standard is based, and the present text draws heavily from them. There have been many contributors to this work both in terms of concepts and in terms of text. This edition clarifies several details and removes the end-of-letter buffer-size adjustments, and redescribes the letter mechanism as a push function.

What transport service does an app need? Data loss Throughput

- some apps (e.g., audio) can tolerate some loss
- other apps (e.g., file transfer, telnet) require 100% reliable data transfer

Timing

 some apps (e.g., Internet telephony, interactive games) require low delay to be "effective"

- some apps (e.g., multimedia) require minimum amount of throughput to be "effective"
- other apps ("elastic apps") make use of whatever throughput they get

Security

• Encryption, data integrity,

Transport service requirements of common apps



Application	Data loss	Throughput	Time Sensitive
file transfer	no loss	elastic	no
e-mail	no loss	elastic	no
Web documents	no loss	elastic	no
real-time audio/video	loss-tolerant	audio: 5kbps-1Mbps	yes, 100's msec
		video:10kbps-5Mbps	
stored audio/video	loss-tolerant	same as above	yes, few secs
interactive games	loss-tolerant	few kbps up	yes, 100's msec
instant messaging	no loss	elastic	yes and no



Internet transport protocols services

TCP service:

- connection-oriented: setup required between client and server processes
- *reliable transport* between sending and receiving process
- flow control: sender won't overwhelm receiver
- congestion control: throttle sender when network overloaded
- *does not provide:* timing, minimum throughput guarantees, security

UDP service:

- unreliable data transfer between sending and receiving process
- does not provide: connection setup, reliability, flow control, congestion control, timing, throughput guarantee, or security

Q: why bother? Why is there a UDP?



Internet apps: application, transport protocols

Application	Application layer protocol	Underlying
Application		transport protocol
e-mail	SMTP [RFC 2821]	ТСР
remote terminal access	Telnet [RFC 854]	ТСР
Web	HTTP [RFC 2616]	ТСР
file transfer	FTP [RFC 959]	ТСР
streaming multimedia	HTTP (eg Youtube),	TCP or UDP
	RTP [RFC 1889]	
Internet telephony	SIP, RTP, proprietary	
	(e.g., Skype)	typically UDP