## ELEN 4017

Network Fundamentals Lecture 21

## Purpose of lecture

Chapter 4: Network Layer

- Recap
- Forwarding
- What's inside a router ?


## Network layer

- transport segment from sending to receiving host
- on sending side encapsulates segments into datagrams
- on rcving side, delivers segments to transport layer
- network layer protocols in every host, router
- router examines header fields in all IP datagrams
 passing through it


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Chapter 4: Network Layer

- Recap
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## Forwarding operation

- Consider a 32 bit destination address.
- If forwarding table had an entry for every possible address, this would require 4 billion entries.
- How can we then reduce this and achieve correct forwarding?


## Ranges of addresses

| Destination Address Range | Link |
| :---: | :---: |
| 11001000000101110001000000000000 through <br> 11001000000101110001011111111111 | 0 |
| ```11001000 00010111 0001100000000000 through 1 1 0 0 1 0 0 0 0 0 0 1 0 1 1 1 0 0 0 1 1 0 0 0 ~ 1 1 1 1 1 1 1 1 ~``` | 1 |
| ```110010000001011100011001 00000000 through 1 1 0 0 1 0 0 0 ~ 0 0 0 1 0 1 1 1 ~ 0 0 0 1 1 1 1 1 ~ 1 1 1 1 1 1 1 1 ~``` | 2 |
| Otherwise | 3 |

## Prefix matching

- For this case, its not necessary to explicitly list each entry in the forwarding table, since we have groups.
- We could employ a prefix match.
- If the key in the prefix table is a prefix of the destination address, then it is matched and then the output interface is used.
- If not, the next key is checked.
- If no match occurs then a default interface is chosen.
- Can you see a problem with the algorithm ?

| Prefix match | Link interface |
| :--- | :---: |
| 110010000001011100010 | 0 |
| 110010000001011100011000 | 1 |
| 110010000001011100011 | 2 |
| otherwise | 3 |

## Longest prefix matching rule

- Since there is a possibility of multiple matches, the rule is extended to match the longest prefix.
- Consider keys for interface 1 and 2.

| Prefix match | Link interface |
| :--- | :---: |
| 110010000001011100010 | 0 |
| 110010000001011100011000 | 1 |
| 110010000001011100011 | 2 |
| Otherwise | 3 |

## An example from telecommunications

- Every subscriber has a unique address - as defined by ISDN numbering plan (E164 standard)
- Consider a Vodacom subscriber number 082123 4567.
- Within SA the number 821234567 is unique
- South Africa also has a country code assigned as 27
- Thus the globally unique number is 27821234567
- Incidentally, the 0 is just a trunk access code, indicating the destination dialled is national.


## Configuring charging bands

- The aim is to configure different charging bands dependent on the destination dialed.
- Thus the following charging plan is required:

| Destination | Price per min |
| :--- | :--- |
| Vodacom subscriber | 1.00 |
| Cell-C subscriber | 1.50 |
| MTN subscriber | 2.00 |
| Destination in U.K | 10.00 |
| Destination in U.S | 15.00 |
| Vodafone partner in U.K | 7.00 |
|  |  |

- This charging can be done by using prefixes.
- The longest prefix matched is the chosen tariff.
- What is the flaw with the method below for SA ?

| Destination | Prefix | Price per min |
| :--- | :--- | :--- |
| Vodacom subscriber | 2782,2772 | 1.00 |
| Cell-C subscriber | 2784 | 1.50 |
| MTN subscriber | 2783 | 2.00 |
| Destination in U.K | 44 | 10.00 |
| Destination in U.S | 1 | 15.00 |
| Vodafone partner in U.K | 44171 | $\mathbf{7 . 0 0}$ |
| Other SA destinations | $\mathbf{2 7}$ | $\mathbf{2 . 5 0}$ |

## Tree indexing for performance

- The longest prefix search is also termed a prefixtree or best fit search.
- The reason for naming it prefix tree is that the data is often indexed in a prefix tree to make searching more efficient.

| Destination | Prefix | Price per min |
| :--- | :--- | :--- |
| Other SA | 27 | 2.50 |
| Vodacom | 2782 | 1.00 |
| Vodacom discounted | 2782123 | 0.80 |
| U.K. | 44 | 10.00 |
| Vodafone partner | 44171 | 7.00 |

