### **Discrete Maths**

#### Data and Information Management: ELEN 3015

School of Electrical and Information Engineering, University of the Witwatersrand

◆□▶ ◆□▶ ◆臣▶ ◆臣▶ 臣 の�?

## Overview

◆□▶ ◆□▶ ◆三▶ ◆三▶ 三三 のへぐ

Discrete maths

Prime Numbers

Greatest Common Divisor (GCD)

**Relative Prime** 

The Euler Totient Function

Modular Arithmetic

Fermat's Little Theorem

◆□▶ ◆□▶ ◆臣▶ ◆臣▶ 臣 の�?

We will look at aspects of number theory that apply to cryptography.

Discrete mathematics is a branch of mathematics that deals with Integers only.

#### 1.1 Prime Numbers

Def - Prime Number: Any integer greater than one that only has 1 and itself as divisors

Prime numbers: 2, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31, 37, 41, 43, 47, 53, 59, 61, 67, 71, 73, 79, 83, 89, 97, 101, 103, 107, 109, 113, ...

Non-prime number is known as a composite

Fundamental theorem of arithmetic: every positive integer (except 1) can be represented in exactly one way as a product of one or more primes (Hardy and Wright 1979, pp. 2-3).

### 1.2 Greatest Common Divisor (GCD)

Largest integer d that divides a and  $b \in \mathbb{Z} \to \mathsf{Greatest}$  Common Divisor of a and b

Notation: d = GCD(a, b)

Example: GCD(12, 16) = 4

Euclidean algorithm can be used to determine GCD

#### 1.3 Relative Prime

Def - Relative prime: When GCD(a, b) = 1, a and  $b \in \mathbb{Z}$ , then a and b are relative prime (also coprime)

In other words, they share no common factors other than 1

Neither a and b need to be prime

Example: GCD(15, 28) = 1, thus 15 and 28 are relative prime

### 1.4 The Euler Totient Function

Def - the totient  $\varphi(n)$  of a positive integer *n* is defined to be the number of positive integers less than *n* that are relative prime to *n*.

$$\varphi(n) = \left\{ egin{array}{cc} n-1, & n \ {
m prime} \ (p-1)(q-1), & n=pq \ {
m with} \ p \ {
m and} \ q \ {
m prime} \end{array} 
ight.$$

For first scenario, note that p (prime) has  $\{1, 2, 3, ..., p-1\}$  as relative primes

Note that the second scenario only represents a small subset of the composite numbers.

◆□▶ ◆□▶ ◆三▶ ◆三▶ 三三 のへぐ

### 1.5 Modular Arithmetic

Discrete maths operates only on integers  $(\mathbb{Z})$ 

Modular arithmetic restricts results to a maximum modulo size

Modulus means remainder after division

### 1.5 Modular Arithmetic

Def - Equivalence / Congruency: Two integers are equivalent under modulus n if their results mod n are equal

Example: 16 mod 7 = 23 mod 7  $\rightarrow$  16  $\equiv$  23 mod 7

### 1.6 Properties of Modular Arithmetic

Modular arithmetic in non-negative integers forms a construct called a commutative ring with the operation + and  $\times.$ 

If every number other than 0 has an inverse under multiplication, the group is called a Galois field. Example: The integers  $a \mod p$  forms a Galois field.

All rings have the properties of associativity and distributivity, commutative rings also have commutativity.

(ロ)、(型)、(E)、(E)、 E) の(の)

### 1.6 Properties of Modular Arithmetic

Example: Modulo 5 Addition

+	0	1	2	3	4
0					
1					
2					
3					
4					

(ロ)、(型)、(E)、(E)、 E) の(の)

### 1.6 Properties of Modular Arithmetic

Example: Modulo 5 Addition

+	0	1	2	3	4
0	0	1	2	3	4
1	1	2	3	4	0
2	2	3	4	0	1
3	3	4	0	1	2
4	4	0	1	2	3

### 1.6 Properties of Modular Arithmetic

Additive identity  $\rightarrow$  *a* + 0 = *a* for any *a*  $\in$   $\mathcal{F}$ 

Additive inverse of an element in  $\mathcal{F}$ :

$$a+(-a)=0$$

Additive inverses:

### 1.6 Properties of Modular Arithmetic

Additive identity  $\rightarrow a + 0 = a$  for any  $a \in \mathcal{F}$ 

Additive inverse of an element in  $\mathcal{F}$ :

$$a+(-a)=0$$

Additive inverses:

- $0 \times 0 \mod 5 = 0$
- 1 × 4 mod 5 = 0
- 2 × 3 mod 5 = 0

◆□▶ ◆□▶ ◆臣▶ ◆臣▶ 臣 の�?

### 1.6 Properties of Modular Arithmetic

Example: Modulo 5 Multiplication

×	0	1	2	3	4
0					
1					
2					
3					
4					

◆□▶ ◆□▶ ◆臣▶ ◆臣▶ 臣 の�?

### 1.6 Properties of Modular Arithmetic

Example: Modulo 5 Multiplication

×	0	1	2	3	4
0	0	0	0	0	0
1	0	1	2	3	4
2	0	2	4	1	3
3	0	3	1	4	2
4	0	4	3	2	1

◆□▶ ◆□▶ ◆臣▶ ◆臣▶ 臣 の�?

### 1.6 Properties of Modular Arithmetic

Multiplicative identity  $\rightarrow a \times e = a$  for any  $a \in \mathcal{F}$ 

e =

### 1.6 Properties of Modular Arithmetic

Multiplicative identity  $\rightarrow a \times e = a$  for any  $a \in \mathcal{F}$ 

e = 1

Multiplicative inverse of an element in  $\mathcal{F}$ :

$$a imes rac{1}{a} = 1$$

Multiplicative Inverses:

### 1.6 Properties of Modular Arithmetic

Multiplicative identity  $\rightarrow a \times e = a$  for any  $a \in \mathcal{F}$ 

e = 1

Multiplicative inverse of an element in  $\mathcal{F}$ :

$$a imes rac{1}{a} = 1$$

Multiplicative Inverses:

- $1 \times 1 \mod 5 = 1$
- 2 × 3 mod 5 = 1
- 4 × 4 mod 5 = 1

◆□▶ ◆□▶ ◆臣▶ ◆臣▶ 臣 の�?

### 1.7 Modulo Inverses

Finite field (Galois Field)  $\rightarrow$  every element except 0 has multiplicative inverse

Ring  $\rightarrow$  not every element might have an inverse

#### 1.7 Modulo Inverses Example: Multiplication Modulo 6

×	0	1	2	3	4	5
0	0	0	0	0	0	0
1	0	1	2	3	4	5
2	0	2	4	0	2	4
3	0	3	0	3	0	3
4	0	4	2	0	4	2
5	0	5	4	3	2	1

2,3 and 4 doesn't have inverses under modulo 6 multiplication 2,3 and 4 not relative prime to 6

Modulo under a prime number  $\rightarrow$  Field (Galois Field), every nonzero element has an inverse

< □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > <

### 1.7 Modulo Inverses Example:

 $4 imes \lambda \equiv 1 \mod 7 o 4\lambda = 7k + 1$ ,  $k \in \mathbb{Z}$ 

General Problem:

 $1 = (\mathsf{a} \times \lambda) \mod n$ 

 $a^{-1} \equiv \lambda \mod n$ 

For Example:  $4^{-1} = 2$  ( $\mathcal{F} = x \mod 5$ )

### 1.8 Fermat's Little Theorem If p is a prime, and a is not a multiple of p:

$$a^{p-1} \equiv 1 \mod p$$

Euler's generalization: if GCD(a, n) = 1:

 $a^{\varphi(n)} \mod n = 1$ 

To compute inverse *x*:

$$x = a^{\varphi(n)-1} \mod n$$

(Can also use Euclid's algorithm)

### 1.9 Properties of Modular Arithmetic

Associativity	[a + (b + c)]mod n = [(a + b) + c]mod n		
	$[a \times (b \times c)]mod n = [(a \times b) \times c] mod n$		
Commutativity	$(a + b) \mod n = (b + a) \mod n$		
	$(a \times b) \mod n = (b \times a) \mod n$		
Distributivity	$(a \times (b + c)) \mod n$		
	= ((a $ imes$ b ) + (a $ imes$ c)) mod n		
Identities	$(a + 0) \mod n = (0 + a) \mod n = a$		
	(a $ imes$ 1) mod n = (1 $ imes$ a) mod n = a		
Inverses	$(a + (-a)) \mod n = 0$		
	$(a imes a^{-1}) m{mod} \ n = 1$		
Reducibility	$(a + b) \mod n$		
	= ((a mod n) + (b mod n)) mod n		
	$(a \times b) \mod n$		
	= ((a mod n) $ imes$ (b mod n)) mod n		

#### 1.10 Euclidean Algorithm

#### Not in the notes!

For any pair of positive integers *a* and *b*, we may find gcd(a, b) by repeated use of division to produce a decreasing sequence of integers  $r_1 > r_2 > \cdots$  as follows.

$$\begin{array}{ll} a = bq_1 + r_1 & 0 < r_1 < b, \\ b = r_1q_2 + r_2 & 0 < r_2 < r_1, \\ r_1 = r_2q_3 + r_3 & 0 < r_3 < r_2, \\ \vdots & \vdots \\ r_{k-3} = r_{k-2}q_{k-1} + r_{k-1} & 0 < r_{k-1} < r_{k-2}, \\ r_{k-2} = r_{k-1}q_k + r_k & 0 < r_k < r_{k-1}, \\ r_{k-1} = r_kq_{k+1} + 0 \end{array}$$

◆□▶ ◆□▶ ◆臣▶ ◆臣▶ 臣 の�?

#### 1.11 Extended Euclidean Algorithm

#### Not in the notes!

For any nonzero integers a and b, there exist integers s and t such that gcd(a, b) = as + bt. Moreover, gcd(a, b) is the smallest positive integer of the form as + bt.

Extended Euclidean Algorithm

$$r_i = r_{i-2} - \left\lfloor \frac{r_{i-2}}{r_{i-1}} \right\rfloor \cdot r_{i-1}$$

< □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > <

### 1.11 Extended Euclidean Algorithm Example: GCD(120,23)

Step	Quotient	Remainder	Expression
1		120	$120 = 120 \times 1 + 23 \times 0$
2		23	$23 = 120 \times 0 + 23 \times 1$
3	5	5	5 = (120  imes 1 + 23  imes 0) - (120  imes 0 + 23  imes 1)  imes 5
			5=120 imes1+23 imes-5

◆□▶ ◆□▶ ◆臣▶ ◆臣▶ 臣 のへぐ

### 1.11 Extended Euclidean Algorithm Example: GCD(120,23)

Step	Quotient	Remainder	Expression
1		120	$120 = 120 \times 1 + 23 \times 0$
2		23	$23 = 120 \times 0 + 23 \times 1$
3	5	5	5 = (120  imes 1 + 23  imes 0) - (120  imes 0 + 23  imes 1)  imes 5
			5=120 imes1+23 imes-5
4	4	3	$3 = 23 - 5 \times 4$
			$3 = (120 \times 0 + 23 \times 1)$ -4(120 -5 $\times$ 23)
			$3 = 120 \times -4 + 23 \times 21$
5	1	2	$2 = 5 - 3 \times 1$
			$2 = (120 \times 1 + 23 \times -5) - (120 \times -4 + 23 \times 21)$
			$2 = 120 \times 5$ - $23 \times 26$
6	1	1	$1 = 3 - 2 \times 1$
			$1 = (120 \times -4 + 23 \times 21) - (120 \times 5 - 23 \times 26)$
			$1 = 120 \times -9 + 23 \times 47$
7	2	0	

### 1.11 Extended Euclidean Algorithm

Extended Euclidean Algorithm can be used to calculate the multiplicative inverse of a number in a ring (if they exist)

From example:  $1 = 120 \times -9 + 23 \times 46$ 

Over the ring mod 120, 23 and 46 are multiplicative inverses of each other

< □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > <

# Summary

< □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > <

Discrete maths

Prime Numbers

Greatest Common Divisor (GCD)

**Relative Prime** 

The Euler Totient Function

Modular Arithmetic

Fermat's Little Theorem

Euclidean Algorithm and Extended Euclidean Algorithm